# Reinforcement Learning on MDPs 

John Langford

Yahoo Research

Backing Material: http://hunch.net/~jl/tutorial/RL.html

## COMS-4771, Columbia

## Reinforcement Learning is Always Relevant



The answer to: "Is this an RL problem?" is always "yes".

The implication: RL theory is broadly applicable.

The other implication: RL theory is often only weakly relevant. (breadth+relevance=hard.)

Understanding a problem as an RL problem is the beginning to solving it. Whenever possible, you want to understand how the problem is special.

## Outline

## 1. Sample Complexity Results

2. Limitations of Sample Complexity

## Markov Decision Process (MDP)

1. $S=$ the number of states in an MDP
2. $A=$ the number of actions/state in an MDP
3. $T=$ the horizon time you care about (or $\gamma=$ discount factor)
4. $O=$ number of observations
5. $\epsilon=$ precision parameter

Important Derived Quantities

$$
V_{t}^{\pi}(s)=E_{(s, a, r)^{t} \sim \pi, \operatorname{MDP}_{s}}\left[\sum_{t^{\prime}=1}^{t} r_{t^{\prime}}\right]
$$

$=$ the value of being in state $s$ and acting according to $\pi$ for $t$ timesteps.

$$
Q_{t}^{\pi}(s, a)=E_{(s, a, r)^{t} \sim \pi, \operatorname{MDP}_{s a}}\left[\sum_{t^{\prime}=1}^{t} r_{t^{\prime}}\right]
$$

$=$ the value of being in state $s$, acting with $a$, and then acting according to $\pi$ for $t$ timesteps.

$$
\pi^{*}(s)=\arg \max _{a} Q_{t}^{\pi^{*}}(s, a)
$$

$=$ recursive definition of optimal policy.

$$
Q_{t}^{*}(s, a)=Q_{t}^{\pi^{*}}(s, a)
$$

$=$ short hand for optimal policy $Q$ values.

## The $E^{3}$ Guarantee

Trace Model $=$ ability to read current state $s$, take action $a$, observe next state $s^{\prime}$ and reward $r$. Notation: $T M: A \rightarrow S \times[0,1]$.

Assume $\operatorname{MDP}\left(S, A, p\left(s^{\prime} \mid s, a\right)\right)$ with horizon $T$

1. Original: +assume mixing time $\tau \Rightarrow \operatorname{Poly}\left(S, A, \tau, \frac{1}{\epsilon}\right)$ samples implies ability to act $\epsilon$ optimal for $T>\tau$.
2. Modified: $\operatorname{Poly}\left(S, A, T, \frac{1}{\epsilon}\right)$ samples implies ability to act $\epsilon \mathrm{op}$ timal for $T$ timesteps.
(2) + mixing assumption implies (1). (2) holds even for deterministic worlds. We'll go through (2).

## $E^{3}$ Theorem Statement

Theorem: There exist an algorithm $E^{3}$ such that for all MDP ( $S, A, T, p\left(s^{\prime} \mid a, s\right)$ ) with rewards $r \in[0,1]$, with probability $1-\delta$, for all except $\operatorname{Poly}\left(S, A, T, \frac{1}{\epsilon}, \ln \frac{1}{\delta}\right)$ steps $Q_{T-t}^{E^{3}} \bmod T\left(s, E^{3}(h)\right) \geq$ $V_{T-t}^{*} \bmod T^{(s)}-\epsilon$ where $h$ is the history of observations.

Suboutline:

1. The Algorithm
2. The Proof

The Known(h) MDP

A state $s$, all actions $a$ leaving $s$ and the probability of their outcomes is known if all actions $a$ leaving $s$ have been executed at least $n$ times.


Initially: known MDP $=$ nothing

## The Known(h) MDP



Then: Reward

Complete dangling action(s) with one state that always has reward 0.

The Known(h) MDP


The Known(h) MDP

(note: the probabilities are empirical counts)

The Unknown (h) MDP

Unknown(h) = Known(h) except the reward is 1 for actions which leave the known states and 0 otherwise.

## Dynamic Program

Fundamental operation: Given MDP $M$ and state $s$,

$$
\mathrm{DP}(M, s, t)=a, v
$$

where $v=$ the maximum expected $T-(t \bmod T)$ reward sum and $a=$ action achieving it.

Computation:

$$
\begin{gathered}
\operatorname{DP}(M, s, t)=\max _{a} E_{s^{\prime}, r \sim M(s, a)^{r}+\operatorname{DP}\left(M, s^{\prime}, t+1\right)}^{\operatorname{DP}(M, s, n T)=0}
\end{gathered}
$$

## $E^{3}(h)$ Explicit Explore or Exploit Algorithm

1. If last $s$ not in Known(h): choose the least previously used action
2. Else:
(a) If $\operatorname{DP}(\operatorname{Unknown}(h))>\epsilon^{\prime}$ then act according $\operatorname{DP}($ Unknown $(h))$ until state is unknown or $t \bmod T=0$ then go to (1).
(b) else act according to DP(Known $(h)$ ).

The proof uses 5(!) MDPs

1. MDP - the true MDP (Imposed by world)
2. $\operatorname{Known}(h)=$ known MDP (Known by $E^{3}$ algorithm)
3. Unknown $(h)=$ unknown MDP (Known by $E^{3}$ algorithm)
4. $\mathrm{MDP}_{\mathrm{K}(\mathrm{h})}=\mathrm{MDP}$ restricted to the known states (exists only in proof)
5. $\operatorname{MDP} \mathrm{M}_{(h)}=\mathrm{MDP}$ restricted to the known states with rewards set to 0 except for escaping rewards. (exists only in proof)

## Proof Sketch:

Simulation Lemma:

$$
\left|\operatorname{DP}\left(\operatorname{MDP}_{\mathrm{U} / \mathrm{K}(h)}\right)-\operatorname{DP}((\mathrm{Un}) \operatorname{Known}(h))\right| \leq \frac{1}{\operatorname{Poly}\left(S, A, T, \ln \frac{1}{\delta}\right)}
$$

Explore/Exploit Lemma:

$$
\mathrm{DP}\left(\mathrm{MDP}_{\mathrm{K}(h)}\right)+T \mathrm{DP}\left(\operatorname{MDP}_{\mathrm{U}(h)}\right) \geq \mathrm{DP}(\mathrm{MDP})
$$

So $n=\operatorname{Poly}\left(S, A, T, \ln \frac{1}{\delta}\right)$ implies ability to simulate on known states to precision $\frac{1}{\operatorname{Poly}\left(S, A, T, \ln \frac{1}{\delta}\right)} \ll \epsilon$. $\Rightarrow$ Explore/Exploit Lemma implies $\operatorname{DP}(\operatorname{MDP})-\operatorname{DP}\left(\operatorname{MDP}_{K(h)}\right)>\epsilon \Rightarrow \operatorname{DP}\left(\operatorname{MDP}_{(h)}\right)>\frac{\epsilon}{T}$ $\Rightarrow$ probability about $\frac{\epsilon}{T}$ of encountering new state if exploring. This can happen only $O\left(\frac{n S A T}{\epsilon}\right)$ times (Using the Chernoff bound). Each exploration uses at most $T$ steps $\Rightarrow$ proof.

## Delayed Q-learning

The theorem can be tightened from $\operatorname{Poly}(S, A)$ to $\widetilde{O}(S A)$ using the Delayed Q-learning algorithm.

## Outline

## 1. Sample Complexity Results

2. Limitations of Sample Complexity

The Limits of Sample Complexity: A lower bound

Theorem: Any algorithm $A$ satisfying the $E^{3}$ statement must use at least $\Omega(T S A)$ actions to explore.
(There are stronger lower bounds, but this is sufficient.)

## Proof

A "Key lock" MDP


States in a chain. One action leads to next state, all the rest lead to the beginning. The final state has an action with reward 1.

## Implications

Lower bound $\Rightarrow$ the really big problems can't be solved.

But the problems are solvable: we solve them every day.
$\Rightarrow$ More or different assumptions are required.

## Related Reading

[ $E^{3}$ ] Michael Kearns and Satinder Singh, "Near Optimal Reinforcement Learning in Polynomial Time", ICML 1998.
[Delayed Q-learning] Alexander Strehl et al, "PAC Model-Free Reinforcement Learning', ICML 2006.
[Sparse Sampling] Michael Kearns, Yishay Mansour, and Andrew Ng, "A Sparse Sampling Algorithm for Near-Optimal Planning in Large Markov Decision Processes", IJCAI 1999.
[Many things] Sham Kakade, On the Sample Complexity of Reinforcement Learning Thesis Gatsby, UCL, 2003.

