### Reinforcement Learning on MDPs

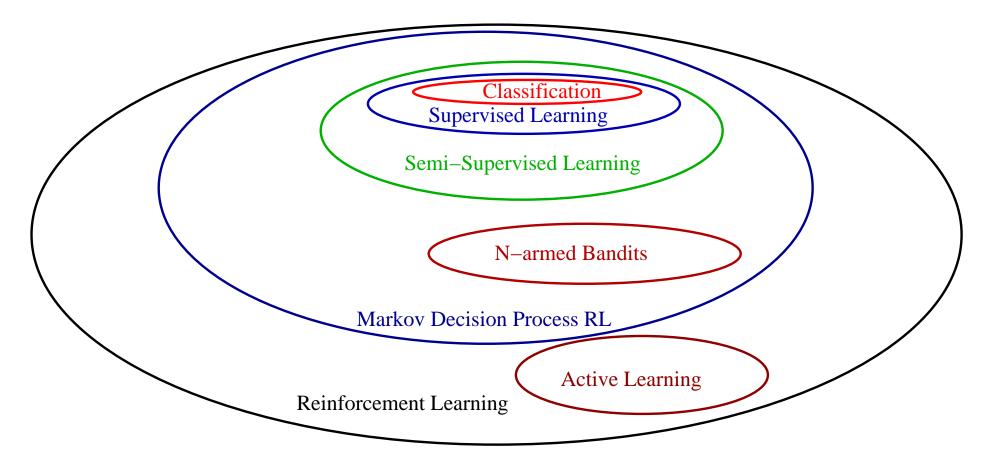
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Backing Material: http://hunch.net/~jl/tutorial/RL.html

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#### Reinforcement Learning is Always Relevant



The answer to: "Is this an RL problem?" is always "yes".

The implication: RL theory is broadly applicable.

The other implication: RL theory is often only weakly relevant. (breadth+relevance=hard.)

Understanding a problem as an RL problem is the *beginning* to solving it. Whenever possible, you want to understand how the problem is special.

# Outline

- 1. Sample Complexity Results
- 2. Limitations of Sample Complexity

Markov Decision Process (MDP)

- 1. S = the number of states in an MDP
- 2. A = the number of actions/state in an MDP
- 3. T = the horizon time you care about (or  $\gamma =$  discount factor)
- 4. O = number of observations
- 5.  $\epsilon$  = precision parameter

Important Derived Quantities

$$V_t^{\pi}(s) = E_{(s,a,r)^t \sim \pi, \mathsf{MDP}_s} \left[ \sum_{t'=1}^t r_{t'} \right]$$

= the value of being in state s and acting according to  $\pi$  for t timesteps.

$$Q_t^{\pi}(s,a) = E_{(s,a,r)^t \sim \pi, \mathsf{MDP}_{sa}} \left[ \sum_{t'=1}^t r_{t'} \right]$$

= the value of being in state s, acting with a, and then acting according to  $\pi$  for t timesteps.

$$\pi^*(s) = \arg\max_a Q_t^{\pi^*}(s,a)$$

= recursive definition of optimal policy.

 $Q_t^*(s,a) = Q_t^{\pi^*}(s,a)$ 

= short hand for optimal policy Q values.

# The $E^3$ Guarantee

Trace Model = ability to read current state s, take action a, observe next state s' and reward r. Notation:  $TM : A \rightarrow S \times [0, 1]$ .

Assume MDP(S, A, p(s'|s, a)) with horizon T

- 1. Original: +assume mixing time  $\tau \Rightarrow \text{Poly}(S, A, \tau, \frac{1}{\epsilon})$  samples implies ability to act  $\epsilon$  optimal for  $T > \tau$ .
- 2. Modified: Poly $(S, A, T, \frac{1}{\epsilon})$  samples implies ability to act  $\epsilon$  optimal for T timesteps.

(2) + mixing assumption implies (1). (2) holds even for deterministic worlds. We'll go through (2).

# $E^3$ Theorem Statement

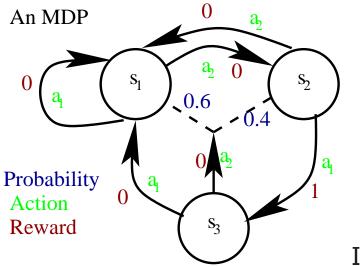
Theorem: There exist an algorithm  $E^3$  such that for all MDP (S, A, T, p(s'|a, s)) with rewards  $r \in [0, 1]$ , with probability  $1 - \delta$ , for all except  $Poly(S, A, T, \frac{1}{\epsilon}, ln \frac{1}{\delta})$  steps  $Q_{T-t \mod T}^{E^3}(s, E^3(h)) \geq V_{T-t \mod T}^*(s) - \epsilon$  where h is the history of observations.

Suboutline:

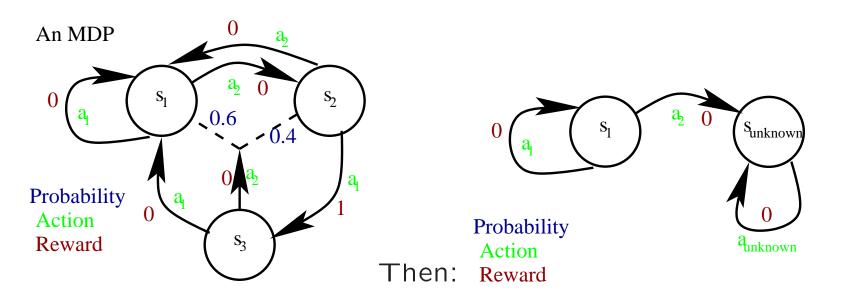
1. The Algorithm

2. The Proof

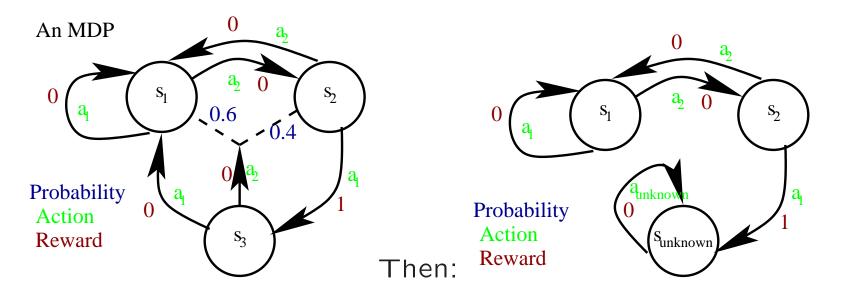
A state s, all actions a leaving s and the probability of their outcomes is known if all actions a leaving s have been executed at least n times.

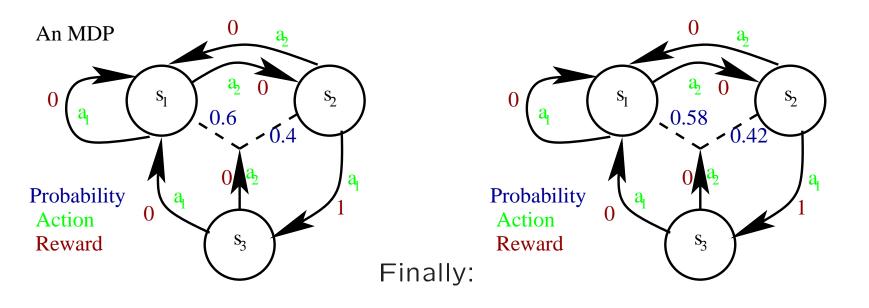


Initially: known MDP = nothing



Complete dangling action(s) with one state that always has reward 0.





(note: the probabilities are empirical counts)

### The Unknown(h) MDP

Unknown(h) = Known(h) except the reward is 1 for actions which leave the known states and 0 otherwise.

#### Dynamic Program

Fundamental operation: Given MDP M and state s,

### $\mathsf{DP}(M, s, t) = a, v$

where v = the maximum expected  $T - (t \mod T)$  reward sum and a = action achieving it.

Computation:

$$DP(M, s, t) = \max_{a} E_{s', r \sim M(s, a)}r + DP(M, s', t + 1)$$
$$DP(M, s, nT) = 0$$

 $E^{3}(h)$  Explicit Explore or Exploit Algorithm

- 1. If last s not in Known(h): choose the least previously used action
- 2. Else:
  - (a) If  $DP(Unknown(h)) > \epsilon'$  then act according DP(Unknown(h))until state is unknown or  $t \mod T = 0$  then go to (1).
  - (b) else act according to DP(Known(h)).

### The proof uses 5(!) MDPs

- 1. MDP the true MDP (Imposed by world)
- 2. Known(h) = known MDP (Known by  $E^3$  algorithm)
- 3. Unknown(h) = unknown MDP (Known by  $E^3$  algorithm)
- 4.  $MDP_{K(h)} = MDP$  restricted to the known states (exists only in proof)
- 5.  $MDP_{U(h)} = MDP$  restricted to the known states with rewards set to 0 except for escaping rewards. (exists only in proof)

#### Proof Sketch:

Simulation Lemma:

 $|\mathsf{DP}(\mathsf{MDP}_{\mathsf{U}/\mathsf{K}(h)}) - \mathsf{DP}((\mathsf{Un})\mathsf{Known}(h))| \leq \frac{1}{\mathsf{Poly}(S, A, T, \ln \frac{1}{\delta})}$ 

Explore/Exploit Lemma:

 $\mathsf{DP}(\mathsf{MDP}_{\mathsf{K}(h)}) + T\mathsf{DP}(\mathsf{MDP}_{\mathsf{U}(h)}) \ge \mathsf{DP}(\mathsf{MDP})$ 

So  $n = \text{Poly}(S, A, T, \ln \frac{1}{\delta})$  implies ability to simulate on known states to precision  $\frac{1}{\text{Poly}(S, A, T, \ln \frac{1}{\delta})} <<\epsilon$ .  $\Rightarrow$  Explore/Exploit Lemma implies  $\text{DP}(\text{MDP}) - \text{DP}(\text{MDP}_{K(h)}) > \epsilon \Rightarrow \text{DP}(\text{MDP}_{U(h)}) > \frac{\epsilon}{T}$  $\Rightarrow$  probability about  $\frac{\epsilon}{T}$  of encountering new state if exploring. This can happen only  $O(\frac{nSAT}{\epsilon})$  times (Using the Chernoff bound). Each exploration uses at most T steps  $\Rightarrow$  proof. Delayed Q-learning

The theorem can be tightened from Poly(S, A) to  $\tilde{O}(SA)$  using the Delayed Q-learning algorithm.

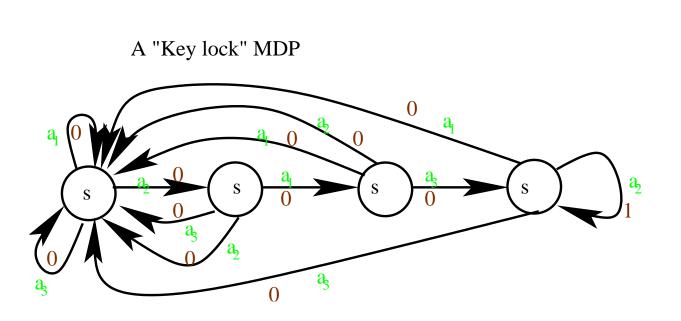
# Outline

- 1. Sample Complexity Results
- 2. Limitations of Sample Complexity

The Limits of Sample Complexity: A lower bound

Theorem: Any algorithm A satisfying the  $E^3$  statement must use at least  $\Omega(TSA)$  actions to explore.

(There are stronger lower bounds, but this is sufficient.)



States in a chain. One action leads to next state, all the rest lead to the beginning. The final state has an action with reward 1.

Proof

#### Implications

Lower bound  $\Rightarrow$  the really big problems can't be solved.

But the problems *are* solvable: we solve them every day.

 $\Rightarrow$ More or different assumptions are required.

#### Related Reading

 $[E^3]$  Michael Kearns and Satinder Singh, "Near Optimal Reinforcement Learning in Polynomial Time", ICML 1998.

[Delayed Q-learning] Alexander Strehl et al, "PAC Model-Free Reinforcement Learning", ICML 2006.

[Sparse Sampling] Michael Kearns, Yishay Mansour, and Andrew Ng, "A Sparse Sampling Algorithm for Near-Optimal Planning in Large Markov Decision Processes", IJCAI 1999.

[Many things] Sham Kakade, On the Sample Complexity of Reinforcement Learning Thesis Gatsby, UCL, 2003.